

Securing Concast

Ken Calvert (calvert@netlab.uky.edu)

Jim Griffioen, Billy Mullins,
Leon Poutievski, Amit Sehgal

ActiveCast Project

Laboratory for Advanced Networking

University of Kentucky, USA

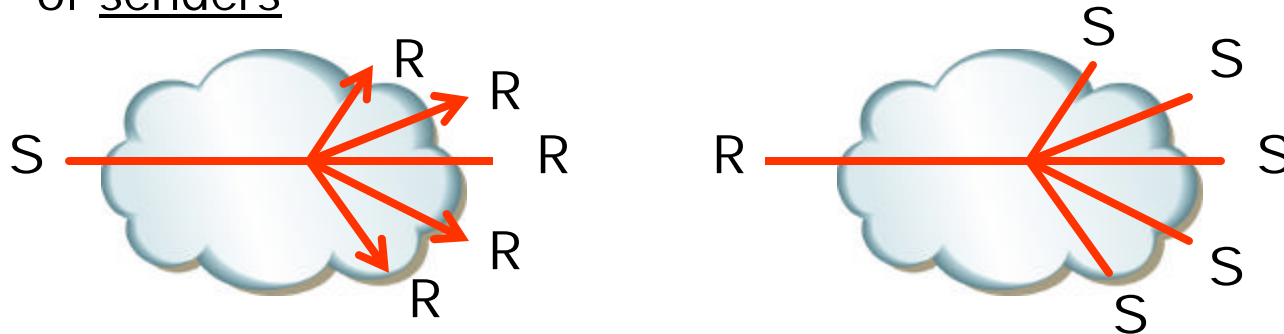
<http://protocols.netlab.uky.edu/~acast/>

Outline

- Concast overview
- GSEC Interest (?)
- Security requirements for concast
- IPsec-based Approach
- Policy considerations

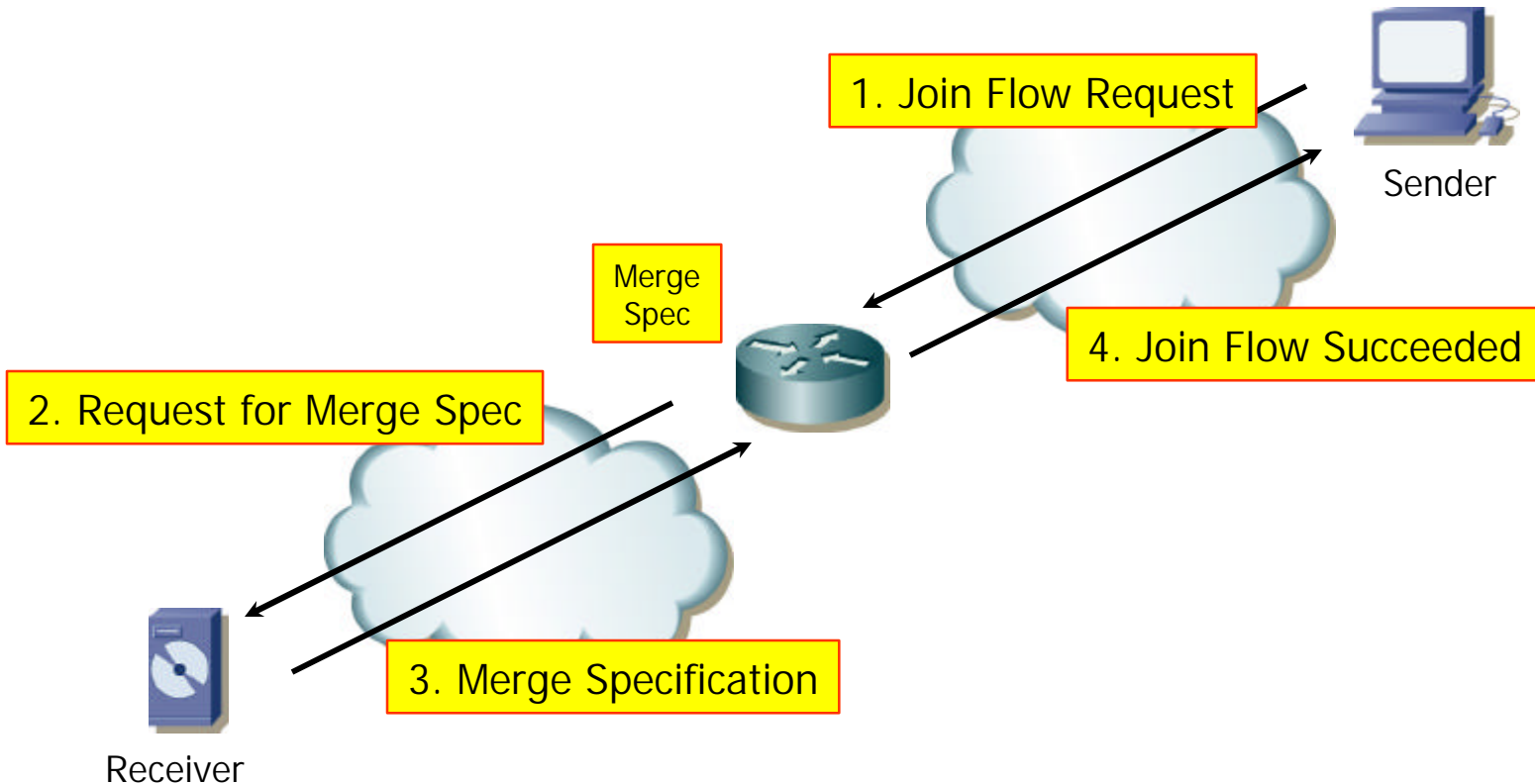
What is Concast?

- Many-to-one network service (inverse multicast)
- Scalability through [abstraction](#)
 - "Inverse Multicast": Single address represents any number of senders

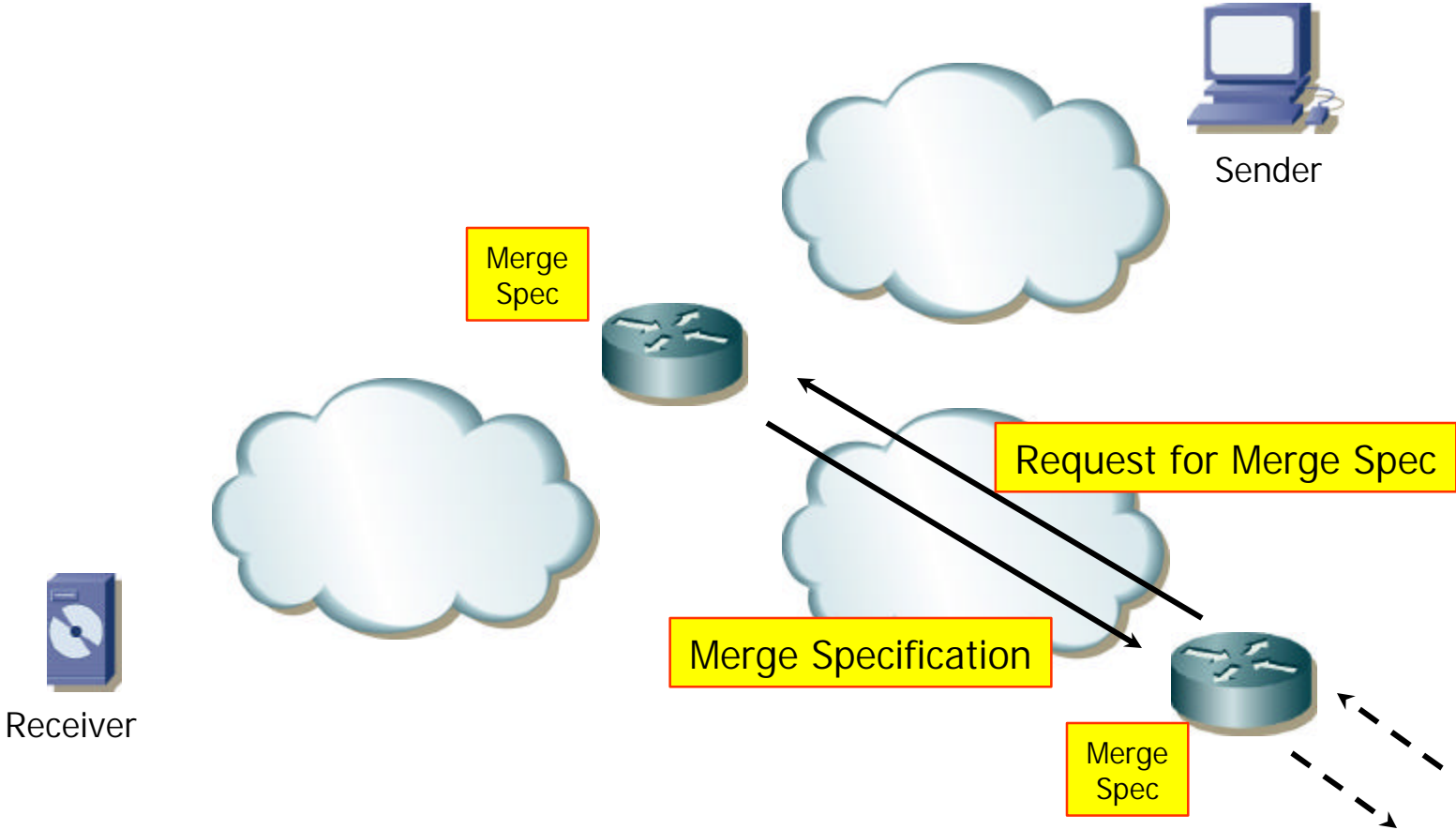


- Network merges messages on route to receiver
...according to a [receiver-supplied merge specification](#)
- Multiple sends result in single message delivery
- [Benefits both receiver and network](#)

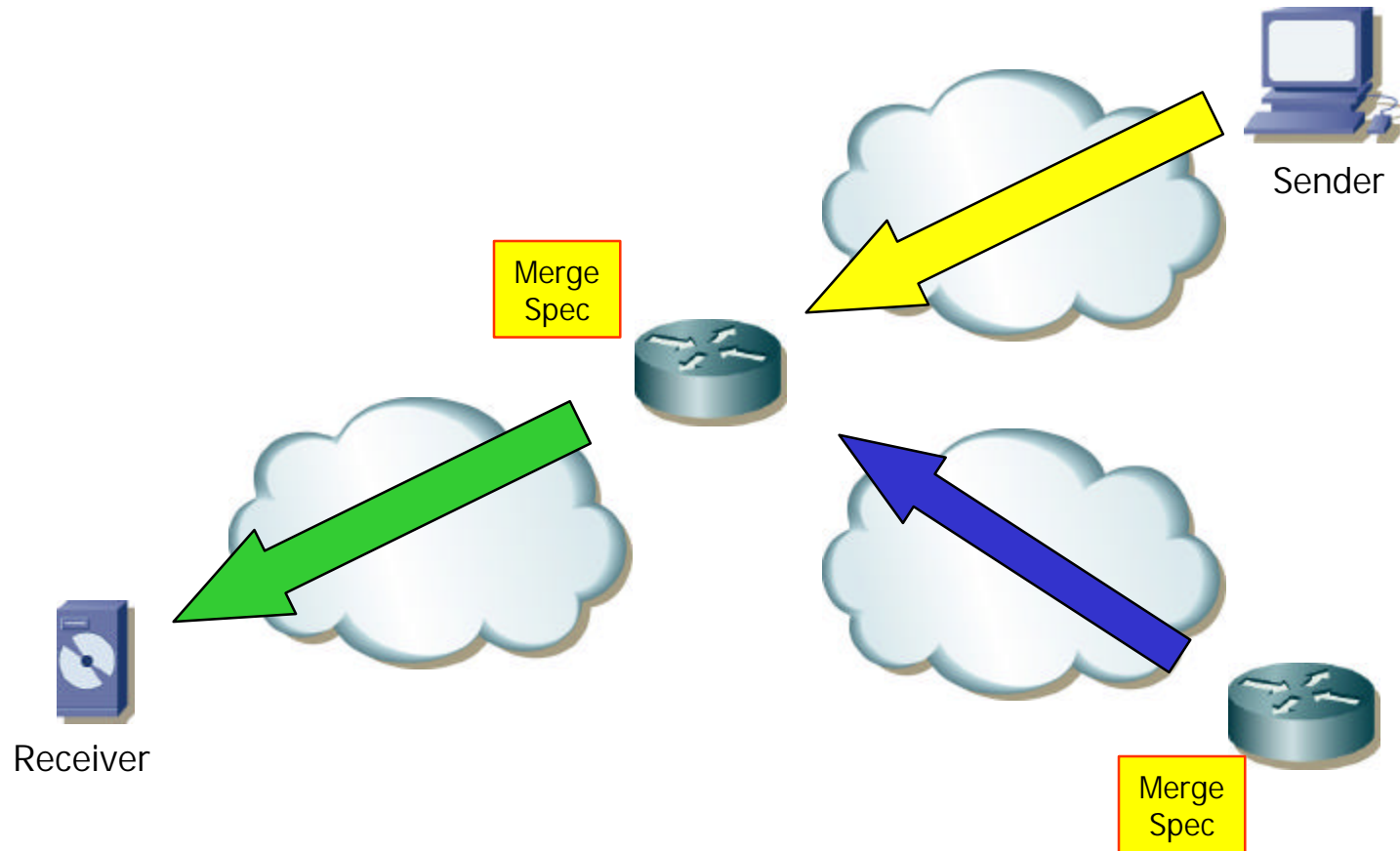
Concast Signaling



Concast Signaling



Concast Data Flow



Why GSEC Might Care (or not)

- A **non-multicast** group application
- A **different** way to secure group feedback
 - With (generic) assistance from the infrastructure
- Somewhat different **policy** issues
 - Third-party policy enforcement (necessary for scaling?)
- **Concast is not currently deployed** (is multicast?)
 - Internet draft (expired) describes the service (no signaling)
- But the ideas and approach may apply in other contexts (e.g. overlays)

Concast Security Requirements: User

- Origin: concast receiver
(which is usually also the multicast sender)
- Requirement: Data integrity, authenticity, confidentiality
- Application-level policy:
 - Which entities can participate as a sender
- Network-level policy:
 - Which **routers** (IP addresses/domains) are trusted to:
 - Perform merging
 - Enforce user policies!

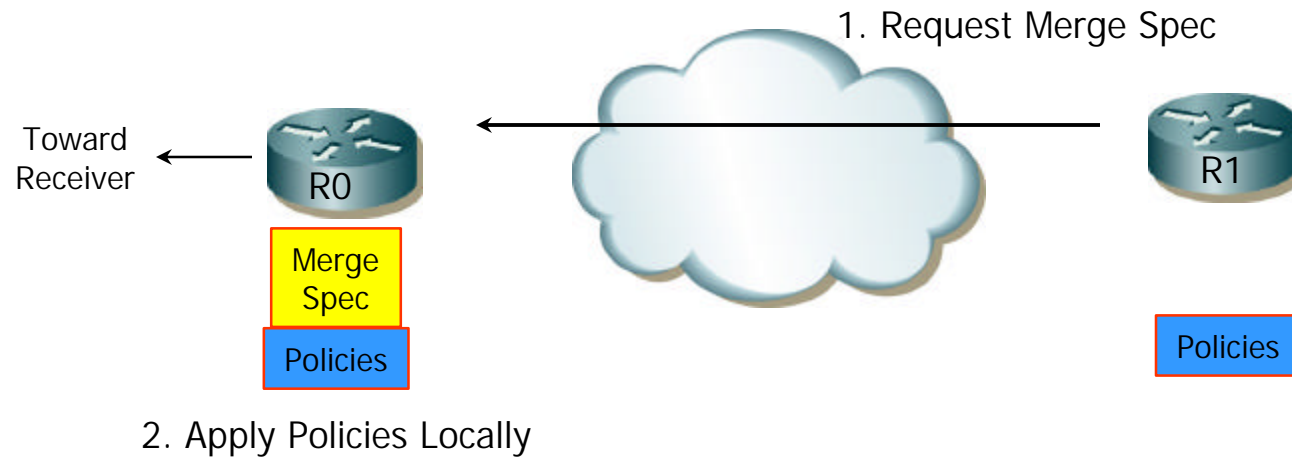
Concast Security Requirements: Provider

- Origin: concast-capable router (or proxy)
- Requirement: access control
 - Ensure that only paying customers use the service
 - Enforce customer policies about session participants
- Network-level policies:
 - Which entities (IP addresses) can participate as senders/receivers
 - Which concast-capable **routers** (IP addresses/domains) can be downstream/upstream of this router
 - Note: policies specify only neighbor relationships

Approach

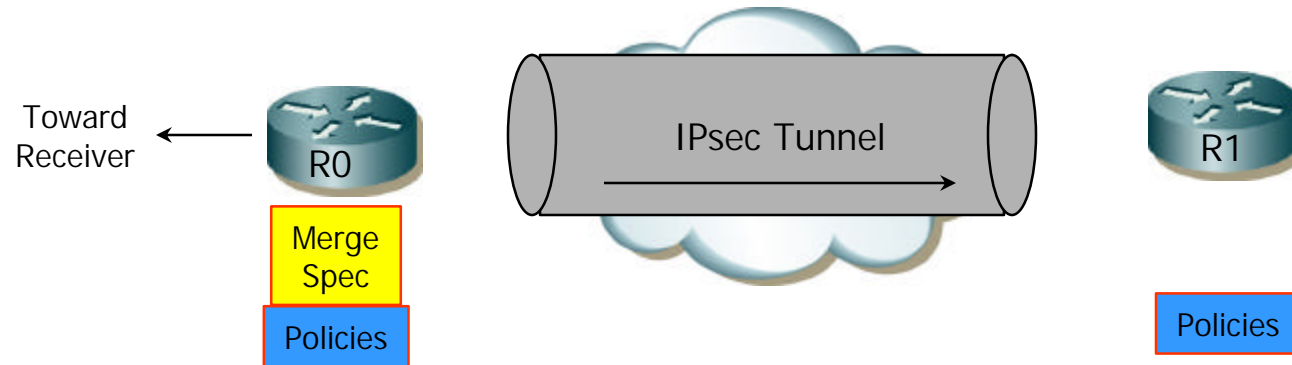
- Idea: secure the signaling; data handled via merge
- Set up IPsec tunnel between concast neighbors at signaling (join/RFM) time
 - No merge specification (including policy) sent in the clear
 - Protocol-specific (CSP) IPsec tunnel
 - Re-use for other flows between same nodes
- Data plane security: up to the application
 - Router trusted to implement user merge spec
 - Add encryption, MAC hooks to the fixed merge framework
 - Applications can instantiate them as desired
 - Transmit keys along with merge specification

Secure Concast Flow Setup



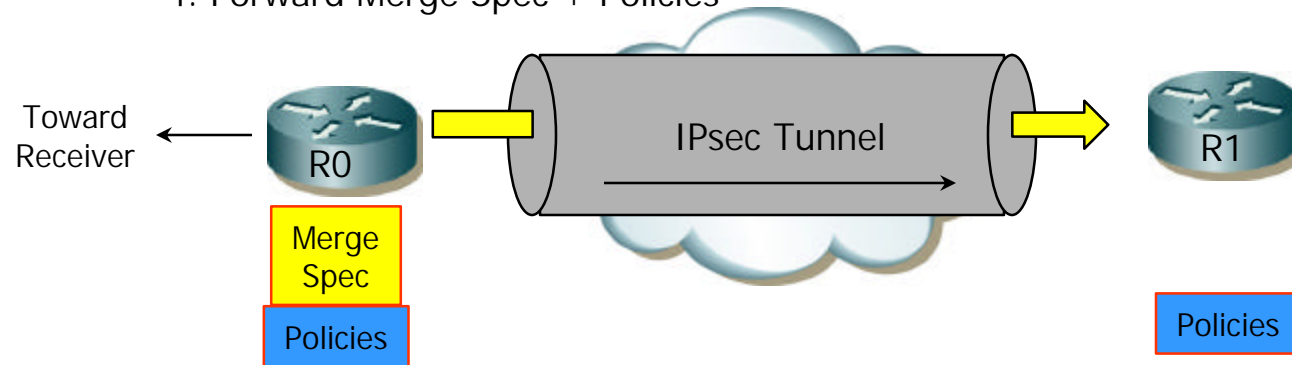
Secure Flow Setup

3. Set up tunnel to upstream neighbor

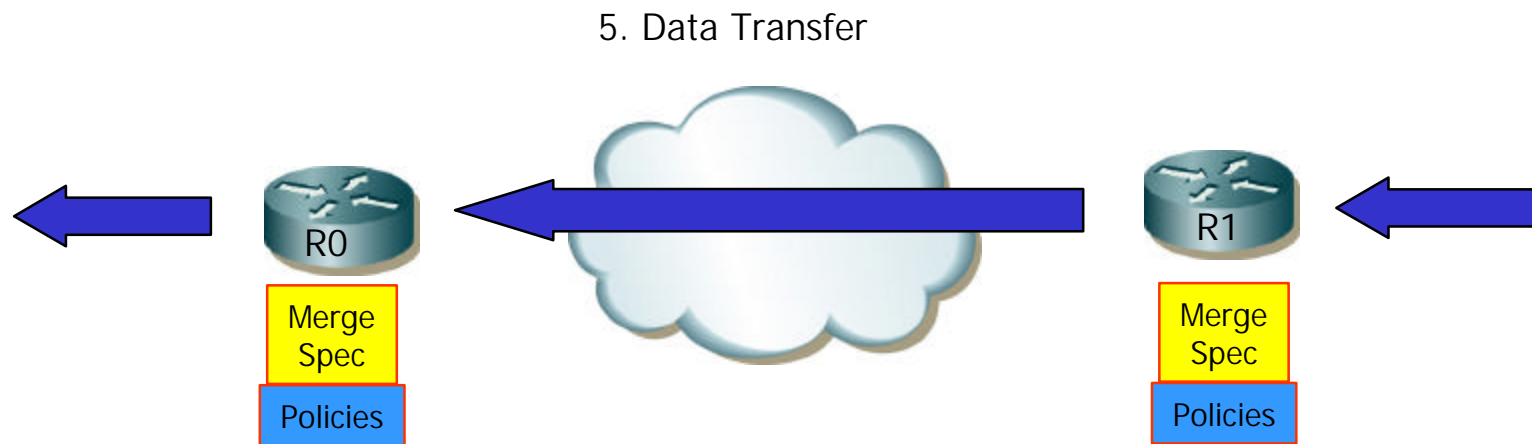


Secure Flow Setup

4. Forward Merge Spec + Policies



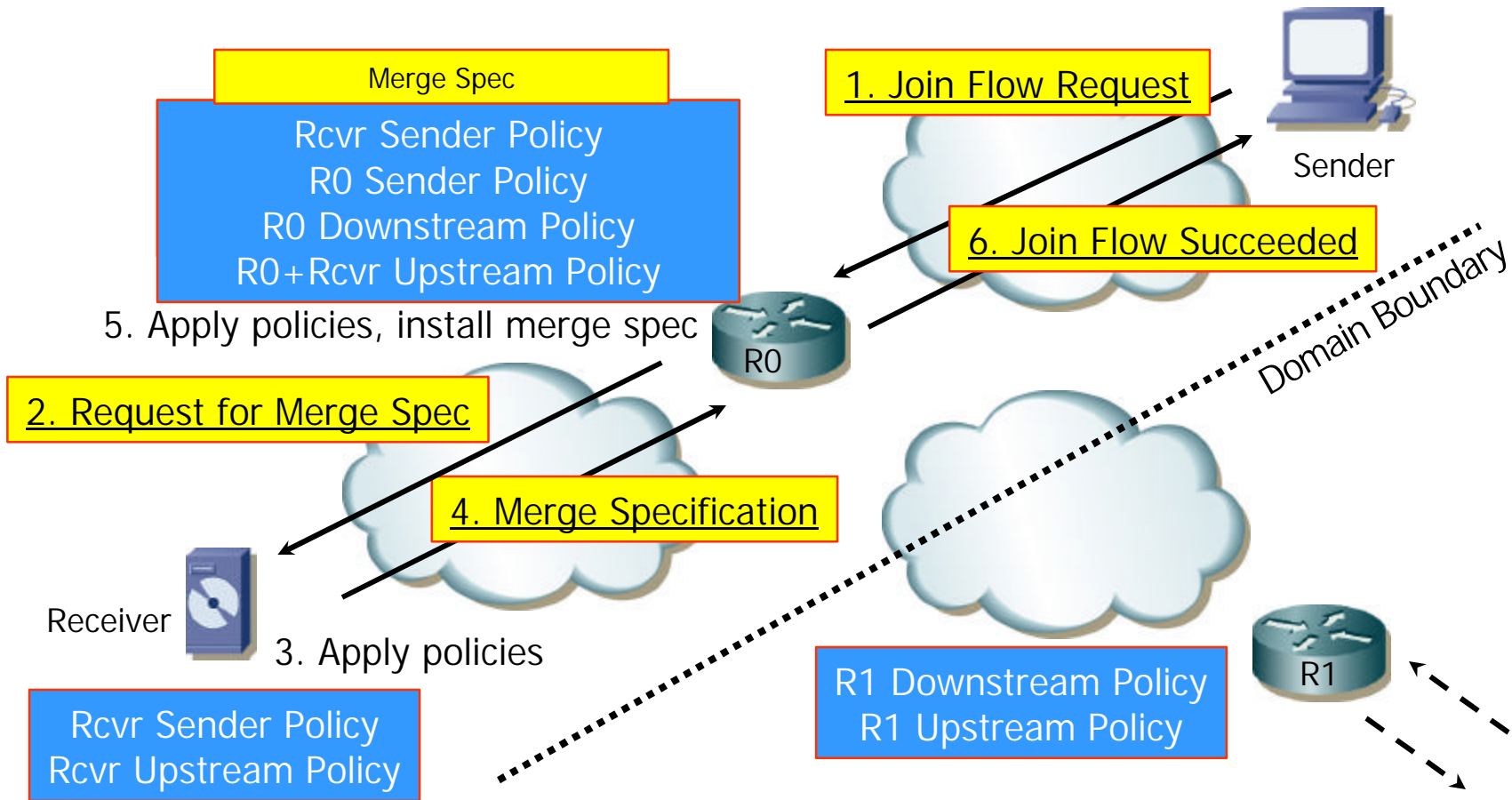
Secure Flow Setup



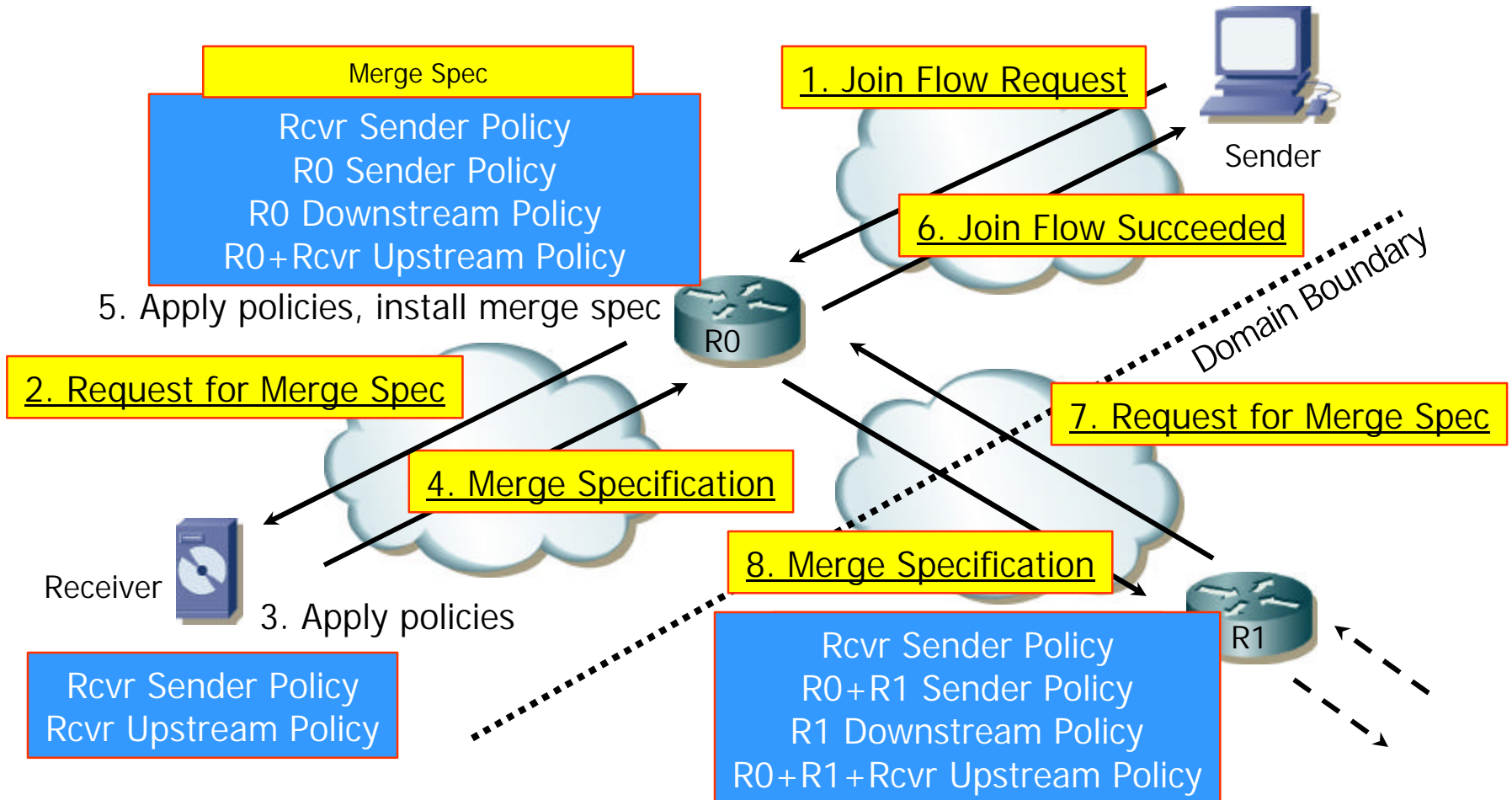
Monotonicity of Policies

- Distributed policy enforcement
- Policies must not be weakened as they propagate upstream from receiver!
 - Need to ensure that **policies at upstream nodes are not weaker than at downstream nodes**
- Access control requirement:
Every principal must satisfy local policy plus policies of all downstream nodes

Policy Monotonicity Requirements



Policy Monotonicity Requirements



Issue: Downstream Neighbor Discovery

- Upstream node does not find out identity of downstream neighbor until merge spec has propagated to that node
 - If downstream node does not satisfy policy, a lot of work may have been wasted!
- Possible solution
 - Upstream node (UN) includes its policy in the setup request message
 - Downstream node (DN) checks itself against UN's policy; if not accepted, DN forwards message without processing it
 - UN trusts DN to forward packets and enforce its policy, but not merge packets (is this sensible?)

Summary

- Concast provides a clean generic framework for doing hop-by-hop app-level processing in the network
- Security can be similarly partitioned:
 - Network is responsible for signaling security and implementing merge specifications accurately
 - Application (merge specification) is responsible for end-to-end security
- This structure helps scalability but raises issues wrt third-party policy enforcement

Questions/comments welcome!

Paper (submission to ICNP) is at:

<http://protocols.netlab.uky.edu/~calvert/private/icnp03-submit.pdf>